

Introduction to Logic [296.617]

Final exam

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1. For each of the following arguments, say if it is valid or not and, in each case, please justify your answer. [20 %]

a)
$$\frac{\begin{array}{l} (1) \text{ If the Moon is made of cheese, then Alan is right.} \\ (2) \text{ The Moon is made of cheese.} \end{array}}{(3) \text{ Alan is right.}}$$

b)
$$\frac{\begin{array}{l} (1) \text{ Every Greek is a man.} \\ (2) \text{ Socrates is a man.} \end{array}}{(3) \text{ Socrates is Greek.}}$$

c)
$$\frac{\begin{array}{l} (1) \ A \rightarrow (B \wedge C) \\ (2) \ A \end{array}}{(3) \ C}$$

2. For each of the following sentences, use the truth-table method to find out whether it is contingent, a tautology, or a logical falsehood. [15 %]

a) $(\varphi \wedge \psi) \rightarrow ((\varphi \vee \chi) \wedge (\psi \vee \chi))$

b) $(\varphi \wedge \neg\psi) \wedge (\varphi \rightarrow \psi)$

c) $(\varphi \vee \neg\psi) \leftrightarrow (\varphi \rightarrow \psi)$

3. Translate the following sentences: [20 %]

a) **From English to the language of propositional logic:** If Mary loves Matthew but Matthew doesn't love Mary, then either Mary marries someone else and is happy or she doesn't and is unhappy.

b) **From the language of first-order logic to English (please make use of the translation key provided):** $((P(x) \rightarrow \neg D(x, y)) \wedge M(x, y)) \rightarrow \neg P(x)$

- $P(x) := x$ is a prime number;
- $D(x, y) := x$ is divisible by y ;
- $M(x, y) := x$ is a multiple of y ;

c) **From English to the language of first-order logic:** Every man that owns a house is happy.

d) **From the language of first-order logic to English (again, please make use of the translation key provided):** $\forall x[[L(x)] \rightarrow \exists y[M(y) \wedge \neg C(x, y)]]$, where

- $L(x) := x$ loves climbing;
- $M(x) := x$ is a mountain;

- $C(x, y) := x$ climbs y .

4. Consider the following sentence: [15 %]

$$\exists x \forall y [\exists z [R(z, z) \rightarrow R(x, y)]]$$

- Is it satisfiable? (provide an example)
 - Is it valid? (if not, provide a counterexample)
5. Prove the following theorem using Natural Deduction Rules for propositional logic (primitive rules only): [15 %]

$$\vdash_{ND} (\varphi \wedge \psi) \rightarrow (\psi \wedge \varphi).$$

6. Prove one of the following sentences using Natural Deduction Rules for first-order logic (primitive rules only): [15 %]
- $\vdash_{ND} \exists x [A(x) \wedge B(x)] \rightarrow \exists x [A(x)];$

Solutions

1. Arguments:

- Valid. The first premise is a conditional, while the second premise states that the antecedent is true. Then it cannot be that both premises are true and the conclusion is false.
- Invalid. There could be the case that Socrates is a man but not Greek.
- Valid. Just like the first one, with the difference that the consequent of the conditional in the first premise is a conjunction. If the whole conjunction is true (as in the argument), then the single components of the conjunction are true.

2. Semantic for propositional logic:

a) Tautology:

φ	ψ	χ	$\varphi \wedge \psi$	$\varphi \vee \chi$	$\psi \vee \chi$	$(\varphi \vee \chi) \wedge (\psi \vee \chi)$	$(\varphi \wedge \psi) \rightarrow ((\varphi \vee \chi) \wedge (\psi \vee \chi))$
1	1	1	1	1	1	1	1
1	1	0	1	1	1	1	1
1	0	1	0	1	1	1	1
1	0	0	0	1	0	0	1
0	1	1	0	1	1	1	1
0	1	0	0	0	1	0	1
0	0	1	0	1	1	1	1
0	0	0	0	0	0	0	1

b) Contradiction

φ	ψ	$\neg\psi$	$\varphi \wedge \neg\psi$	$\varphi \rightarrow \psi$	$(\varphi \wedge \neg\psi) \wedge (\varphi \rightarrow \psi)$
1	1	0	0	1	0
1	0	1	1	0	0
0	1	0	0	1	0
0	0	1	0	1	0

c) Contingent

φ	ψ	$\neg\psi$	$\varphi \vee \neg\psi$	$\varphi \rightarrow \psi$	$(\varphi \vee \neg\psi) \leftrightarrow (\varphi \rightarrow \psi)$
1	1	0	1	1	1
1	0	1	1	0	0
0	1	0	0	1	0
0	0	1	1	1	1

3. Translations:

- a) $(L(x, y) \wedge \neg L(y, x)) \rightarrow (M(x, z) \wedge H(x)) \vee (\neg M(x, z) \wedge \neg H(x));$
 i. $L(x, y) := x$ loves y ;
 ii. $M(x, y) := x$ marries y ;
 iii. $H(x) := x$ is happy;
- b) If x is a prime number then x is not divisible. But x is a multiple of y . Then x is not a prime number.
- c) $\forall x[(M(x) \wedge H(x)) \rightarrow Hp(x)].$
 i. $M(x) := x$ is a man;
 ii. $H(x) := x$ is a house;
 iii. $Hp(x) := x$ is happy;
- d) For every man that loves climbing there is a mountain that he cannot climb.

4. First order semantics:

- a) $\exists x \forall y [\exists z [R(z, z) \rightarrow R(x, y)]]$ is satisfiable by the following structure $\langle \mathfrak{M}, s \rangle$:
- Domain: $\mathcal{D} = \{1, 2, 3\}$;
 - $R^{\mathfrak{M}}(n, m) := n \leq m$, for all $n, m \in \mathcal{D}$;
 - $s(x) = 1, s(y) = 2, s(z) = 3$.
- b) On the other hand, $\exists x \forall y [\exists z [R(z, z) \rightarrow R(x, y)]]$ is not valid. Consider as a counterexample the same structure as before but with a different variable assignment:
- Domain: $\mathcal{D} = \{1, 2, 3\}$;
 - $R^{\mathfrak{M}}(n, m) := n \leq m$, for all $n, m \in \mathcal{D}$;
 - $s'(x) = 3, s'(y) = 2, s'(z) = 1$.

5. Natural Deduction for propositional logic:

$$\frac{\frac{\frac{[\varphi \wedge \psi]^1}{\psi} \langle \wedge E \rangle \quad \frac{[\varphi \wedge \psi]^1}{\varphi} \langle \wedge E \rangle}{\psi \wedge \varphi} \langle \wedge I \rangle}{(\varphi \wedge \psi) \rightarrow (\psi \wedge \varphi)} \langle \rightarrow I \rangle$$

6. Natural Deduction for first order logic:

a)

$$\frac{\frac{\frac{[A(y) \wedge B(y)]^2}{A(y)} \langle \wedge E \rangle}{\exists x[A(x)]} \langle \exists I \rangle}{\frac{[\exists x[A(x) \wedge B(x)]]^1}{\exists x[A(x)]} \langle \exists E \rangle} \langle \rightarrow I \rangle$$